# CSE 539: Applied Cryptography Lec 3: Secret Sharing

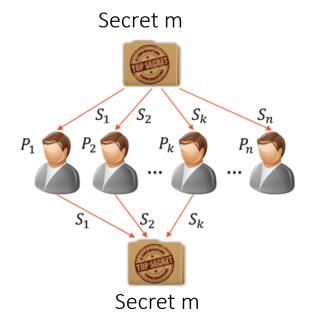
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## Outline

- General Secret Sharing and Examples
- Applications of Secret Sharing
- Constructions
  - "One-time Pad"?
  - Threshold Secret Sharing (Shamir, Blakely 1979)
- Issues

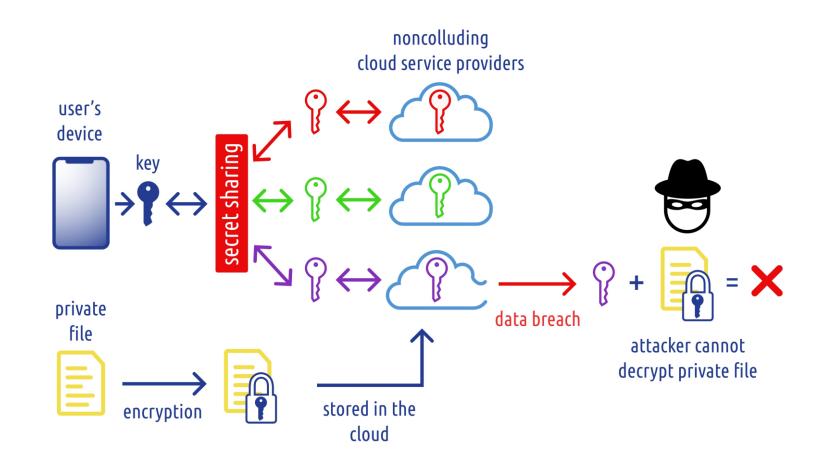
## General Secret Sharing

- m Secret to be shared
- P Set of participants
- => A qualified subsets of can reconstruct m



- Formally, secret sharing scheme allows share a secret m among n parties such that for a fixed number t < n, the following conditions are satisfied.</li>
  - If < t parties get together, then they get no additional information about the secret.
  - If > t parties get together, then they can correctly reconstruct the secret

### Applications of Secret Sharing



### Applications of Secret Sharing

Secure and Efficient Metering [Naor and Pinkas, Eurocrypt 1998] shares Audit Agency **Client Machines Proof of k visits** share **Reconstruct** secret

## Applications of Secret Sharing

- Threshold Signature Sharing
- RSA Signatures
- ...

## Outline

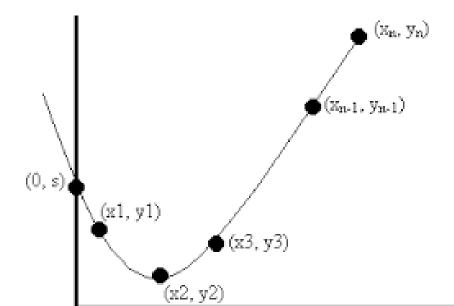
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#### **Basic Solution of Secret Sharing**

• OTP?

t-out-of-n secret sharing

Share(*m*):  $f_1, \dots, f_{t-1} \leftarrow \mathbb{Z}_p$  $f(\mathbf{x}) \coloneqq m + \sum_{j=1}^{t-1} f_j \mathbf{x}^j$ for i = 1 to n:  $s_i \coloneqq (i, f(i) \% p)$  $\mathcal{M} = \mathbb{Z}_p$ p : prime return  $\mathbf{s} = (s_1, \ldots, s_n)$ n < pReconstruct({ $s_i | i \in U$ }):  $t \leq n$  $f(\mathbf{x}) \coloneqq$  unique degree-(t - 1)polynomial mod *p* passing through points  $\{s_i \mid i \in U\}$ return f(0)



https://joyofcryptography.com/pdf/chap3.pdf

• Example: Suppose the secret being shared is m = 7, construct 3-out-of-5 secret sharing over  $Z_{11}$ .

- Sample quiz: Suppose that the secret is m = 7. Which of the following shares are valid in 3-out-of-5 secret sharing scheme
- □(1,1), (2,8), (5,8) □(1,1), (2,8), (3,5)
- **(**1,1), (2,8), (6,9)
- **(**2,8), (3,6), (4,6)

- n users have shared two secrets using Shamir secret sharing.
  - User *i* has a share  $s_i = (i, y_i)$  of the secret *m*, and a share  $s_i' = (i, y_i')$  of the secret *m'*. Both sets of shares use the same prime modulus p, and have the same threshold
- Suppose each user *i* locally computes  $z_i = (y_i + y'_i)\% p$
- Are the resulting  $\{(i, z_i), i \leq n\}$  a valid secret-sharing of the secret m+m'?

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#### Issues

- Honest dealer assumed
- Verifiable Secret Sharing schemes tolerate a faulty dealer
  - Security is computational