# PRACTICAL PRIVACY-PRESERVING K-MEANS CLUSTERING

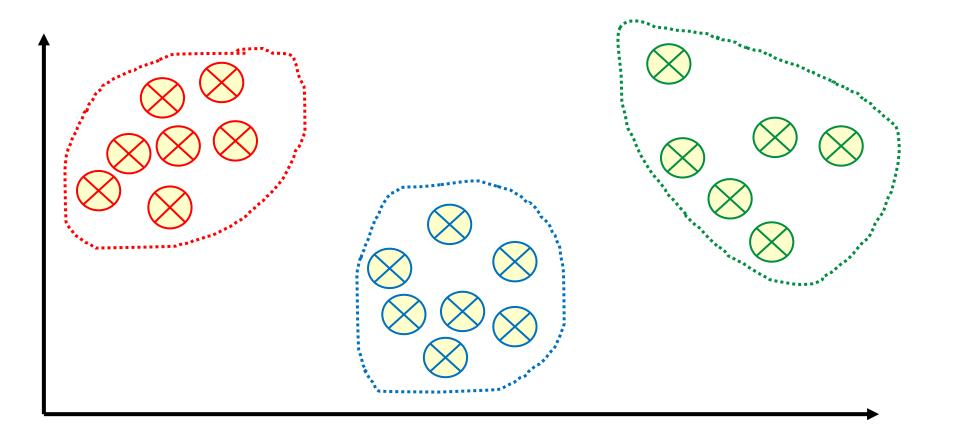
Recording: https://www.youtube.com/watch?v=g7292mN-yAI

Payman Mohassel (Facebook) Mike Rosulek (Oregon State University) **Ni Trieu** (UC Berkeley)



## WHY DO DRIVINGY DRIGHRANG CLUSTERING?

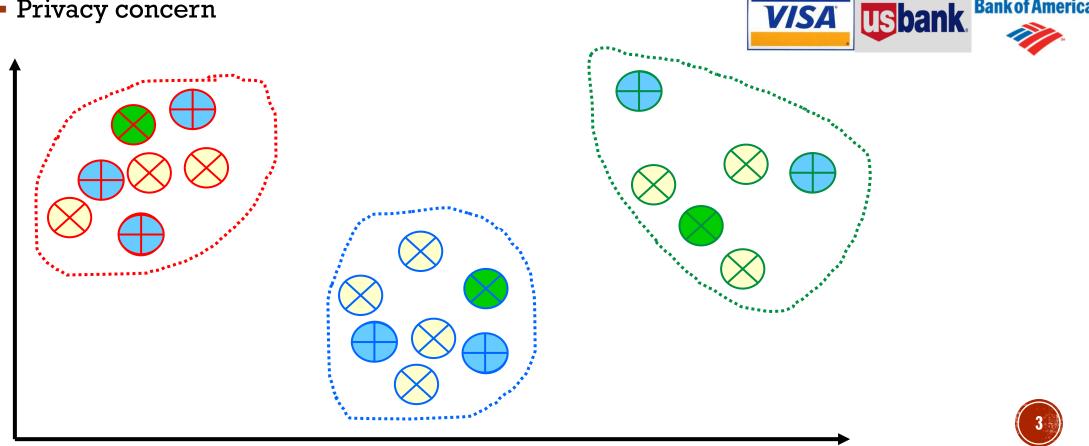
- What is clustering?
  - The process of grouping a set of objects into classes of similar objects



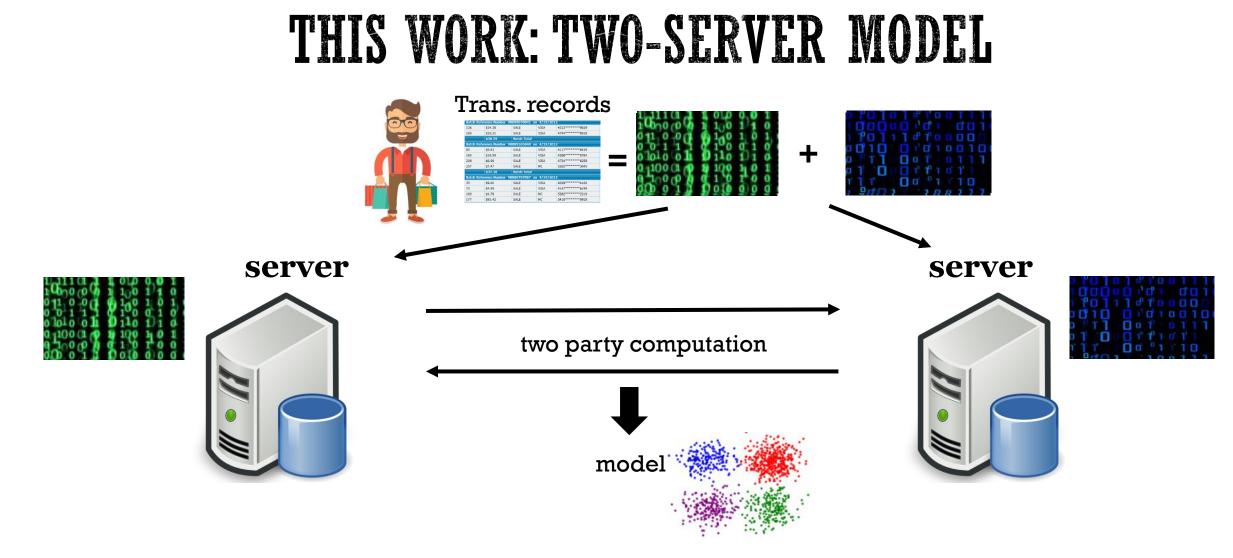


# WHY DO PRIVACY PRESERVING CLUSTERING?

- What is clustering?
  - The process of grouping a set of objects into classes of similar objects
- Why is privacy preserving clustering important?
  - Data comes from different sources: transaction, genome
  - Privacy concern



**Bank of America** 



- Our scheme:
  - More efficient than previous work
  - Scale to large datasets

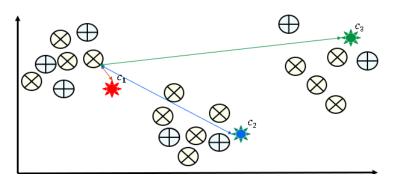


## CLUSTERING ALGORITHMS

• K-mean clustering O(nmt), where dataset size n, number clusters m, number iterations t

This work





 $m \ll n$  and  $t \ll n$ 

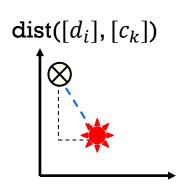
- K-mean clustering is faster than hierarchical clustering
- K-mean clustering is widely used in many applications

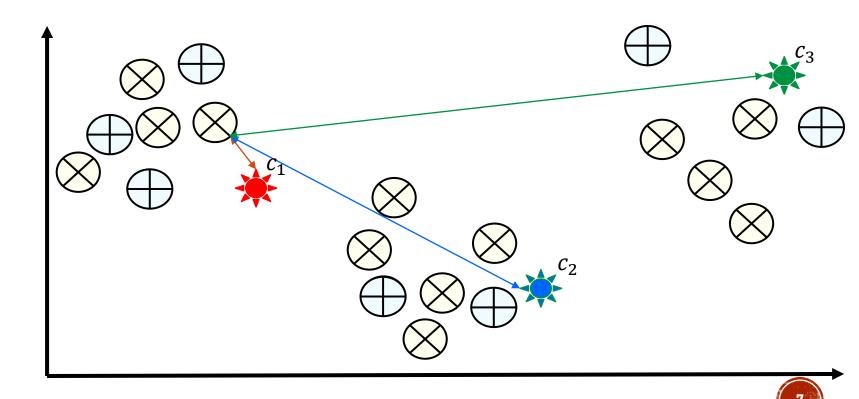


- Privacy-preserving K-mean Clustering has received less attention than other supervised ML problems
- Previous work [...,XHYCZ17, JA18]
  - Not provide a full privacy guarantee (e.g. reveal the intermediate cluster centers)
  - and/or heavy on public key operations
  - $\Rightarrow$  inefficient when the dataset is large.
- Our work:
  - Full privacy guarantee
  - Based on symmetric key operations
    - Efficient secure Euclidean distance
    - Customized minimum circuit on shared values
    - K-means clustering
  - $\Rightarrow$  5 orders of magnitude faster than prior work



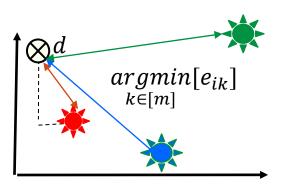
- Inputs:
  - Secret shared dataset as  $\{[d_1], ..., [d_n]\}$
  - #clusters m
- Algorithm:
  - Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
  - Until clustering converges (or other stopping criterion):
    - For each data point d<sub>i</sub>:
      - For each cluster c<sub>k</sub>:
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])

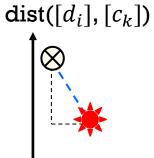


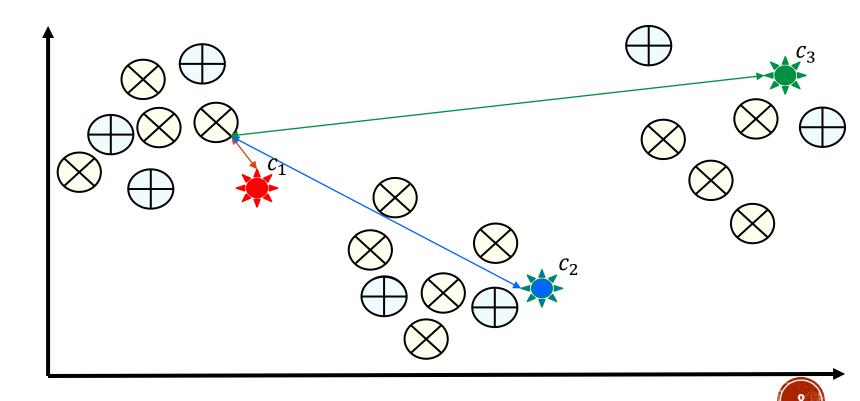


- Inputs:
  - Secret shared dataset as  $\{[d_1], \dots, [d_n]\}$
  - #clusters m
- Algorithm:
  - Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
  - Until clustering converges (or other stopping criterion):
    - For each data point d<sub>i</sub>:
      - For each cluster  $c_k$ :
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
      - Find closest cluster center

 $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$ 



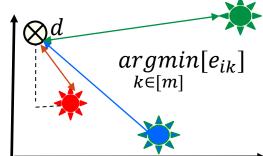


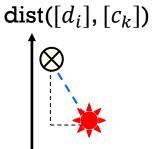


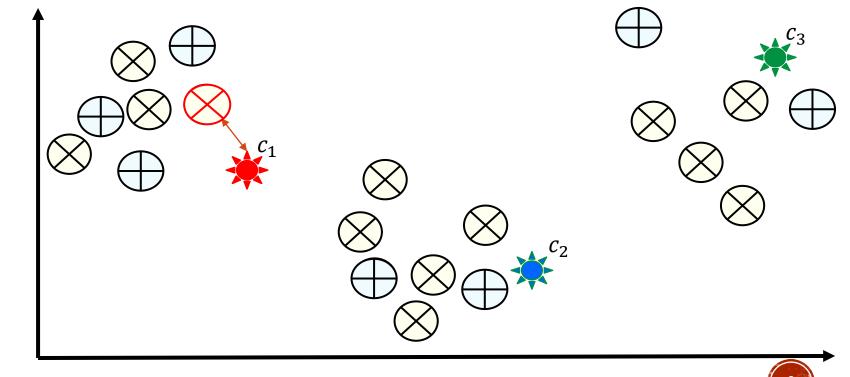
- Inputs:
  - Secret share dataset as  $\{[d_1], ..., [d_n]\}$
  - #clusters m
- Algorithm:
  - Pick m random clusters  $\{[c_1], ..., [c_m]\}$
  - Until clustering converges (or other stopping criterion):
    - For each data point  $d_i$ :
      - For each cluster  $c_k$ :
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
      - Find closest cluster center

 $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$ 

Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>





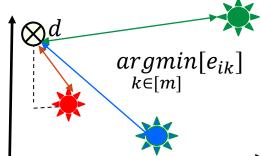


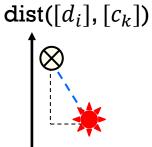
#### Inputs:

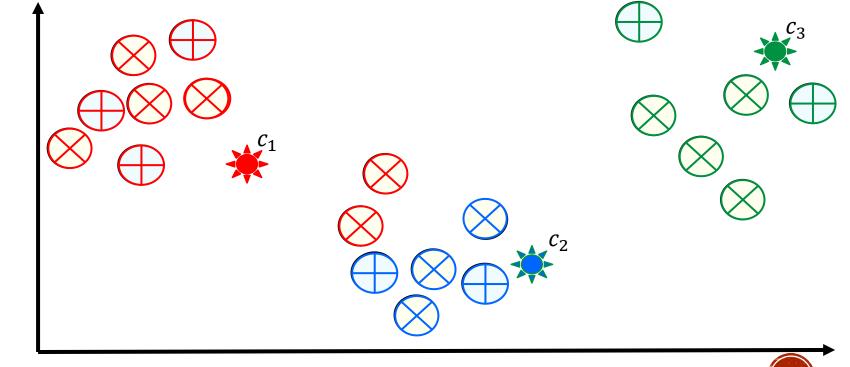
- Secret share dataset as  $\{[d_1], \dots, [d_n]\}$
- #clusters m
- Algorithm:
  - Pick m random clusters  $\{[c_1], ..., [c_m]\}$
  - Until clustering converges (or other stopping criterion):
    - For each data point d<sub>i</sub>:
      - For each cluster  $c_k$ :
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
      - Find closest cluster center

 $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$ 

Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>







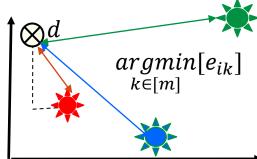
#### Inputs:

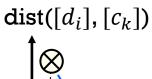
- Secret share dataset as  $\{[d_1], \dots, [d_n]\}$
- #clusters m
- Algorithm:
  - Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
  - Until clustering converges (or other stopping criterion):
    - For each data point d<sub>i</sub>:
      - For each cluster  $c_k$ :
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
      - Find closest cluster center

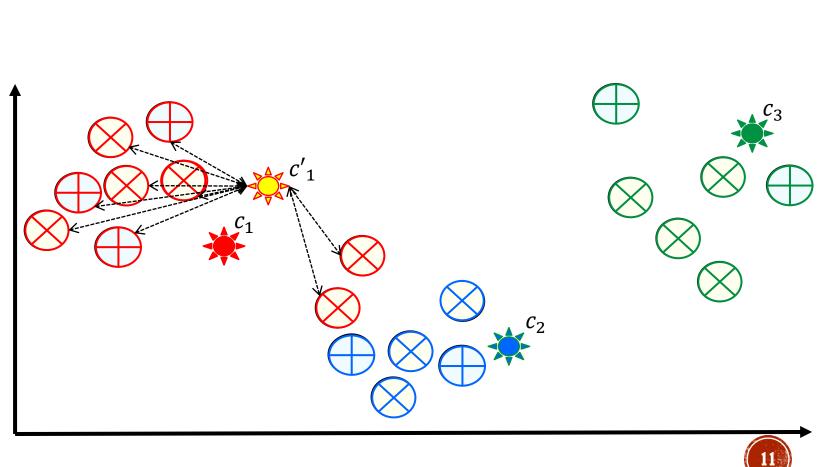
 $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$ 

- Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>
- For each cluster  $c_k$ :
  - Compute the average of the points assigned to the cluster as

 $[c'_k] = avg([d_i]), d_i \in C_k$ 







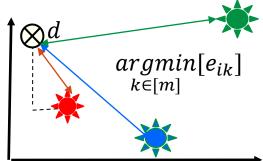
- Inputs:
  - Secret share dataset as  $\{[d_1], ..., [d_n]\}$
  - #clusters m
- Algorithm:
  - Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
  - Until clustering *converges* (or other stopping criterion):
    - For each data point d<sub>i</sub>:
      - For each cluster  $c_k$ :
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
      - Find closest cluster center

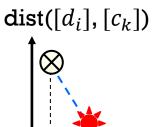
 $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$ 

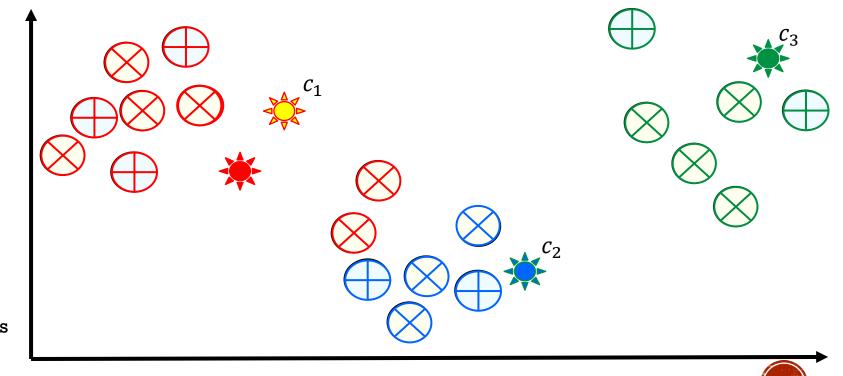
- Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>
- For each cluster  $c_k$ :
  - Compute the average of the points assigned to the cluster as

 $[c'_k] = avg([d_i]), d_i \in C_k$ 

• Update  $c_k = c'_k$ 







#### Inputs:

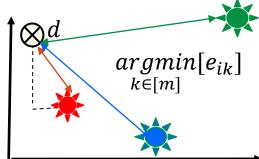
- Secret share dataset as  $\{[d_1], \dots, [d_n]\}$
- #clusters m
- Algorithm:
  - Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
  - Until clustering converges (or other stopping criterion):
    - For each data point  $d_i$ :
      - For each cluster c<sub>k</sub>:
        - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
      - Find closest cluster center

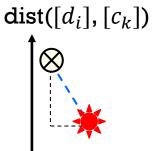
 $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$ 

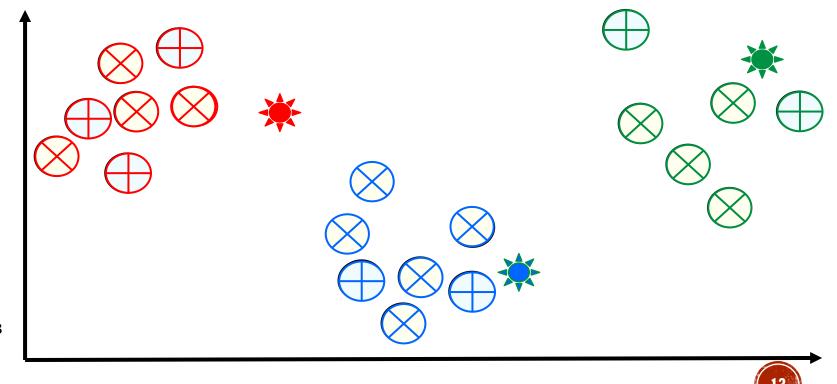
- Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>
- For each cluster  $c_k$ :
  - Compute the average of the points assigned to the cluster as

 $[c'_k] = avg([d_i]), d_i \in C_k$ 

• Update  $c_k = c'_k$ 







#### Inputs:

- Secret share dataset as  $\{[d_1], \dots, [d_n]\}$
- #clusters m
- Algorithm:

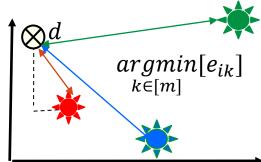
S

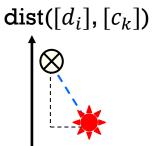
S

- Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
- Until clustering *converges* (or other stopping criterion):
  - For each data point d<sub>i</sub>:
    - For each cluster  $c_k$ :
      - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
    - Find closest cluster center
      - $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$
    - Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>
  - For each cluster  $c_k$ :
    - Compute the average of the points assigned to the cluster as

 $[c'_k] = avg([d_i]), d_i \in C_k$ 

• Update  $c_k = c'_k$ 

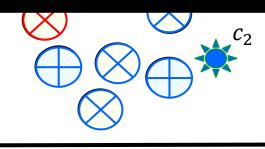




### **Questions:**

 How to securely compute Euclidian distance?
 How to find k\* and assign data to cluster without revealing which cluster?

3. How to obliviously update the cluster?





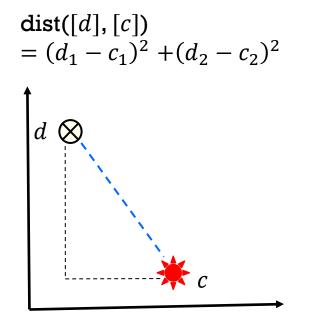






# SECURE EUCLIDIAN DISTANCE





- Squared Euclidian distance dist([d], [c]) =  $\sum_{i=1}^{\ell} (d_i c_i)^2$ , where:
  - d,c is  $\ell$  dimensional vector
  - Each party holds arithmetic share  $d_i = d_i^A + d_i^B$ ,  $c_i = c_i^A + c_i^B$

$$\begin{aligned} \mathsf{dist}([d], [c]) &= \sum_{i=1}^{\ell} \left( d_i^A + d_i^B - c_i^A - c_i^B \right)^2 \\ &= \sum_{i=1}^{\ell} \left( \left( d_i^A - c_i^A \right) + \left( d_i^B - c_i^B \right) \right)^2 \\ &= \sum_{i=1}^{\ell} \left( d_i^A - c_i^A \right)^2 + \sum_{i=1}^{\ell} \left( d_i^B - c_i^B \right)^2 + 2\sum_{i=1}^{\ell} \left( d_i^A - c_i^A \right) \left( d_i^B - c_i^B \right) \end{aligned}$$

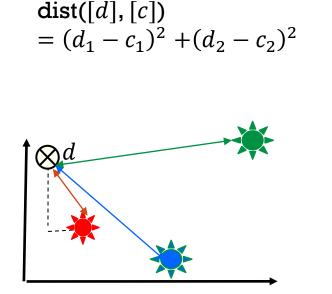
- Alice locally computes  $\sum_{i=1}^{\ell} (d_i^A c_i^A)^2$
- Bob locally computes  $\sum_{i=1}^{\ell} (d_i^B c_i^B)^2$
- Parties jointly compute inner product  $\sum_{i=1}^{\ell} (d_i^A c_i^A) (d_i^B c_i^B)$





# SECURE EUCLIDIAN DISTANCE





- Parties jointly compute inner product  $\sum_{i=1}^{\ell} (d_i^A c_i^A) (d_i^B c_i^B)$ 
  - Previous work: using generic inner product
  - Our observation:
    - Need to compute inner products between 1 point d vs many clusters c
    - Point  $d_i$  is fix during all iterations
    - #cluster << #datapoint</p>
  - We rewrite

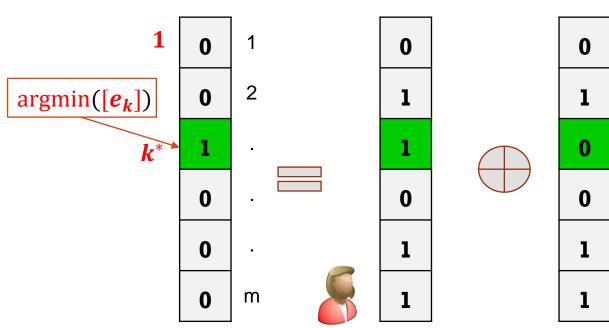
 $(d_{i}^{A} - c_{i}^{A})(d_{i}^{B} - c_{i}^{B}) = d_{i}^{A} (d_{i}^{B} - c_{i}^{B}) - d_{i}^{B} c_{i}^{A} + c_{i}^{B} c_{i}^{A}$ 

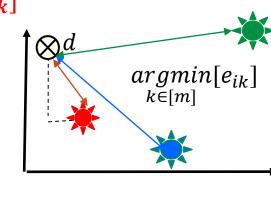
- Applying inner product based on Oblivious Transfer (OT)
  - Reuse OT for  $d_i^A$  when computing  $d_i^A (d_i^B c_i^B)$
  - Reuse OT for  $d_i^B$  when computing  $d_i^B c_i^A$
- => 32-148x faster than prior work



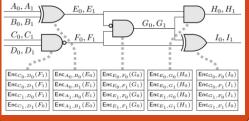
## ASSIGN DATA TO CLUSTER

- Input: shared Euclidian distances between a point and m clusters  $[e_k]$
- Goal:
  - Find closest cluster  $[k^*] = \underset{k \in [m]}{\operatorname{argmin}}[e_k]$
  - Assign the point [d] to the closest cluster  $[c_R]$
- Most of prior works reveal  $k^*$
- Ours: hide k<sup>\*</sup>
  - Presenting  $k^*$  as a binary vector





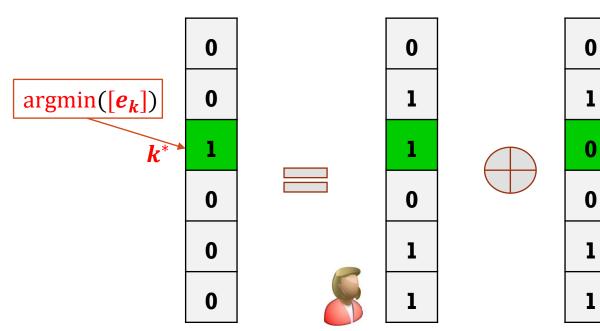
Garbled Circuit



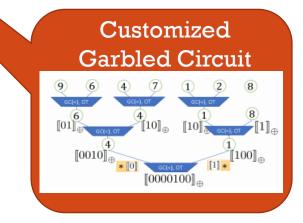


## ASSIGN DATA TO CLUSTER

- Input: shared Euclidian distances between a point and m clusters  $[e_k]$
- Goal:
  - Find shortest distance  $[k^*] = \underset{k \in [m]}{\operatorname{argmin}}[e_k]$
  - Assign the point [d] to the closest cluster  $[c_k]$
- Most of prior works reveal  $k^*$
- Ours: hide k<sup>\*</sup>
  - Present ing  $k^*$  as a binary vector



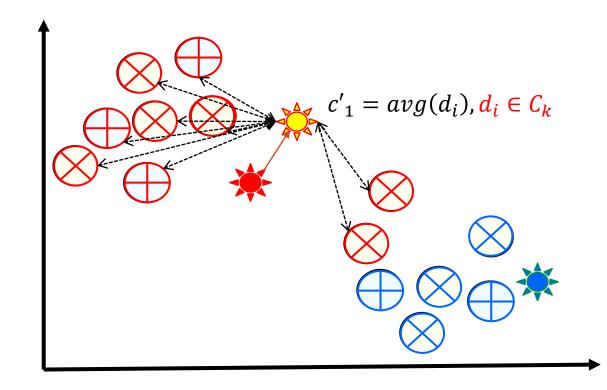
- $[e_k]$ : arithmetic shares
- Garble circuit: Yao shares
- $\Rightarrow$  convert arithmetic shares to Yao's
- $\Rightarrow$  Not cheap!
- $\Rightarrow$  We design a better circuit (see paper)





## UPDATE CLUSTER

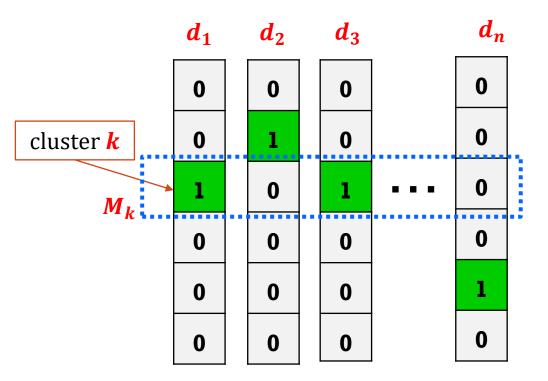
• Recalculate the new cluster center by  $c_k = a\nu g(d_i), d_i \in C_k$ 





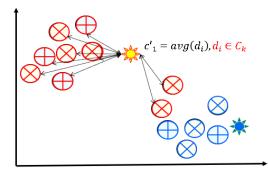
## UPDATE CLUSTER

- Recalculate the new cluster center by  $c_k = a\nu g(d_i), d_i \in C_k$
- From previous step, we have a bit vector indicated whether  $d_i \in C_k$
- Average  $avg(d_i)$  can be computed as:



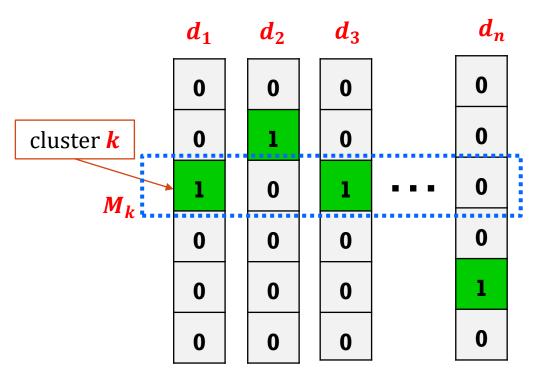
$$c_k = \frac{\sum_{i=1}^n M_k[i] * d_i}{\sum_{i=1}^n M_k[i]}$$

where  $M_k$  and  $d_i$  are secretly shared among two parties



# UPDATE CLUSTER

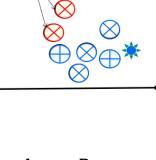
- Recalculate the new cluster center by  $c_k = avg(d_i), d_i \in C_k$
- From previous step, we have a bit vector indicated whether  $d_i \in C_k$
- Average  $avg(d_i)$  can be computed as:



$$c_{k} = \frac{\sum_{i=1}^{n} M_{k}[i] * d_{i}}{\sum_{i=1}^{n} M_{k}[i]} = \frac{\sum_{i=1}^{n} (M_{k}^{A}[i] \oplus M_{k}^{B}[i]) * (d_{i}^{A} + d_{i}^{B})}{\sum_{i=1}^{n} (M_{k}^{A}[i] \oplus M_{k}^{B}[i])}$$

where  $M_k$  and  $d_i$  are secretly shared among two parties

- Direct Solution:
  - Convert Boolean share to arithmetic share
  - Secure multiplication
- Better Solution based Oblivious Transfer (see paper)



 $c'_1 = avg(d_i), d_i \in C_k$ 

#### Inputs:

- Secret share dataset as  $\{[d_1], \dots, [d_n]\}$
- #clusters m
- Algorithm:

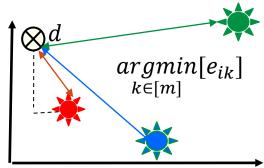
S

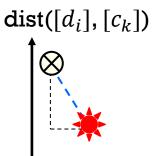
S

- Pick m random clusters {[c<sub>1</sub>], ..., [c<sub>m</sub>]}
- Until clustering converges (or other stopping criterion):
  - For each data point d<sub>i</sub>:
    - For each cluster  $c_k$ :
      - Compute the distance between d<sub>i</sub> and c<sub>k</sub> as [e<sub>ik</sub>] =dist([d<sub>i</sub>], [c<sub>k</sub>])
    - Find closest cluster center
      - $[k^*] = \underset{k \in [m]}{argmin}[e_{ik}]$
    - Assign the data point d<sub>i</sub> to the closest cluster center k<sup>\*</sup>
  - For each cluster c<sub>k</sub>:
    - Compute the average of the points assigned to the cluster as

 $[c'_k] = avg([d_i]), d_i \in C_k$ 

• Update  $c_k = c'_k$ 





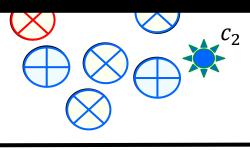
### **Questions:**

 1. How to securely compute Euclidian distance?

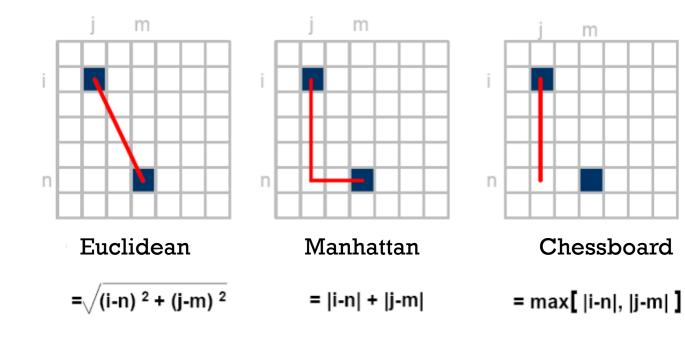
 2. How to find k\* and assign data to cluster

 without revealing which cluster?

3. How to obliviously-update the cluster?



# PERFORMANCE: DISTANCE MEASURES



 Manhattan metric and Chessboard metric are considered as alternative distance metrics in some ML applications.

Our amortized Secure Euclidean Squared Distance (SESD) cost is 8.9 × - 38.5× faster than the cost of computing a Manhattan distance, and 10.5 × - 40.5× faster than that of Chessboard distance.

	Distance	Dimension d					
	Metric	2	3	4	10		
	Manhattan	1.163	1.623	1.96	4.763		
	Chessboard	1.222	1.711	2.294	5.791		
	$\{K = 4, T = 10\}$	0.094	0.155	0.219	0.474		
SESD	$\{K = 4, T = 20\}$	0.079	0.123	0.164	0.398		
SESD	$\{K = 16, T = 10\}$	0 <mark>.036</mark>	0.043	0.066	0.172		
	$\{K = 16, T = 20\}$	0.031	0.042	0.063	0.163		

Running time in millisecond per a distance metric with d dimension. Data size is  $n = 2^{12}$ , K, T is the number of clusters, and number of iterations, respectively.

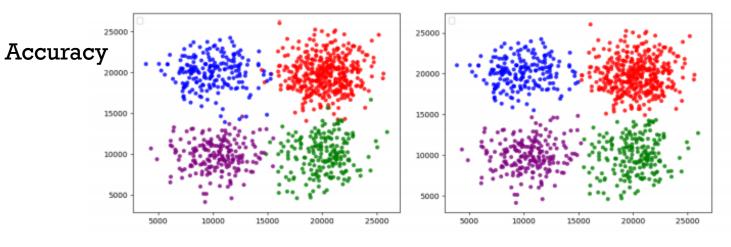


# PERFORMANCE

### **Experiments with Generated Dataset**

Parameters		RunTime (minute)				Communication (MB)					
	n	K	Т	Distance (SESD)	Assign Points to Clusters	Update Centroids	Total	Distance (SESD)	Assign Points to Clusters	Update Centroids	Total
-	104	0	10	0.65	1.14	0.13	1.92	200	2330	10	2540
		2	20	0.95	2.29	0.26	3.5	398	4660	20	4878
		5	10	0.73	4.61	0.47	5.81	496	8760	40	9296
-			20	1.18	9.23	0.94	11.35	989	17520	80	18589
	10 <sup>5</sup>	2	10	5.69	11.12	1.2	18.02	1932	21400	140	23472
			20	10.38	22.25	2.4	35.04	3985	42800	280	47065
		5	10	5.77	47.18	5.13	58.09	4969	85630	340	90939
		<b>1</b>	20	11.13	94.35	10.27	115.75	9927	171260	680	181867

Running time in minute and communication cost of our privacy-preserving clustering protocol, where n, K is the size of database and the number of clusters, respectively, T is number of iterations, dimension d = 2, and bit-length = 3

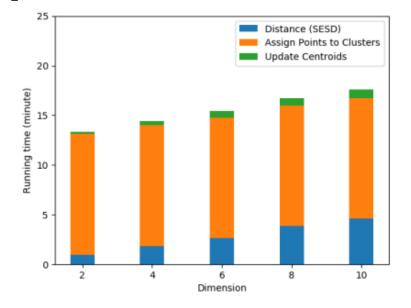


#### (a) Ground Truth Model [2]

(b) Plaintext and Privacy-Preserving K-means Model.

Comparison of accuracy for privacy-preserving, plain-text, and ground truth model. Our privacy-preserving model achieves the same accuracy as the plain-text model, which reaches 95% accuracy compared to the expected ideal clusters

#### **Experiments with Different Dimensions**



Running time (in minute) of our privacy-preserving clustering protocol, where dataset size is 10,000 points, dimension is  $d \in \{2, 4, 6, 8, 10\}$ , the number of cluster and iterations are 9 and 15, respectively

24

# THANK YOU

- Paper: https://eprint.iacr.org/2019/1158.pdf
- <u>Code: https://github.com/osu-crypto/secure-kmean-clustering</u>