# CSE 539: Applied Cryptography Week 13: Oblivious Transfer

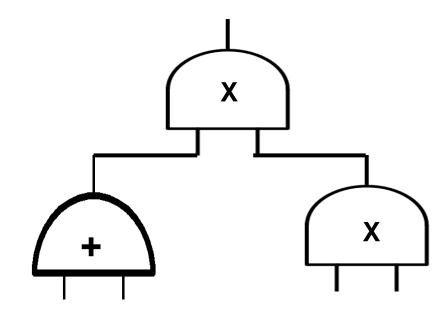
Ni Trieu (ASU)

#### Reading:

- <a href="https://web.engr.oregonstate.edu/~rosulekm/cryptabit/1-overview.pdf">https://web.engr.oregonstate.edu/~rosulekm/cryptabit/1-overview.pdf</a>
- https://securecomputation.org/docs/ch1-introduction.pdf
- <a href="https://securecomputation.org/docs/ch2-definingmpc.pdf">https://securecomputation.org/docs/ch2-definingmpc.pdf</a>

#### Recap: Secure Computation

- Secure computation is a magic box
  - Yao's Protocol (Garbled Circuit)



# Recap: Yao's Protocol

x > y?

- Input domain:  $x, y \in \{1,2\}$ 
  - Alice's input: x = 1
  - Bob's input: y = 2
- Strawman solution:
  - Alice does the following:
    - Write truth table of function f(x, y) = x > y?
    - For each possible input, choose random cryptographic key
    - Encrypt each output with corresponding keys
    - Randomly permute ciphertexts, send to Bob

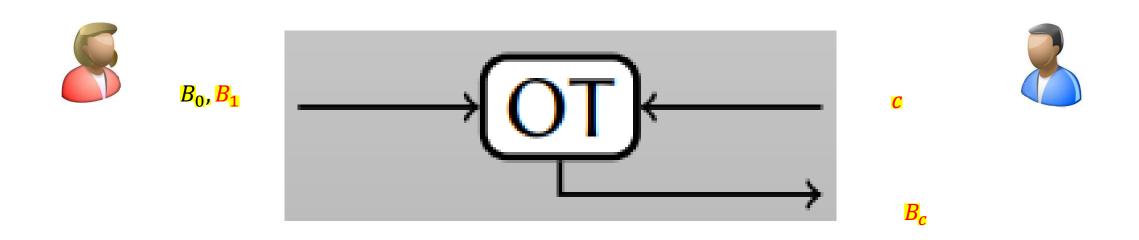
Somehow Bob obtains only "correct" encrypted keys:  $A_x$ ,  $B_y$ 

- Bob learns only f(x, y)
- [Informal security]:
  - Alice learns nothing
  - Bob learns  $A_x$ , but not x. Bob can only decrypt  $Enc_{A_x,B_y}(f(x,y))$



x	у	f(x,y)
$A_1$	$B_1$	$Enc_{A_1,B_1}(f(1,1))$
$A_2$	$B_3$	$Enc_{A_2,B_3}(f(2,3))$
$A_1$	$B_3$	$Enc_{A_1,B_3}(f(1,3))$
$A_2$	$B_1$	$Enc_{A_2,B_1}(f(2,1))$
$A_2$	$B_2$	$Enc_{A_2,B_2}(f(2,2))$
$A_1$	$B_2$	$Enc_{A_1,B_2}(f(1,2))$
$A_3$	$B_1$	$Enc_{A_3,B_1}(f(3,1))$
$A_3$	$B_2$	$Enc_{A_3,B_2}(f(3,2))$
$A_3$	$B_3$	$Enc_{A_3,B_3}(f(3,3))$

# Oblivious Transfer Functionality



Oblivious Transfer (OT) refers to the setting where a sender with two input strings  $(m_0, m_1)$  interacts with a receiver who has an input choice bit b. As the result, the receiver learns mb without learning anything about  $m_{1-b}$ , while the sender learns nothing about b.

## Oblivious Transfer Construction





# Oblivious Transfer: Security





#### OBLIVIOUS TRANSFER EXTENSION

[Beaver'96, Ishai-Kilian-Nissim-Petrank'03]

- OT (using PK) is expensive
- Few OTs+ symmetric keys => many OTs [Ishai-Kilian-Nissim-Petrank'03]
  - But still need to communicate  $O(\kappa)$  bits per random OT, where  $\kappa$  is security parameter



### Oblivious Transfer Construction (reading)

- https://www.iacr.org/archive/crypto2003/27290145/27290145.pdf
- https://eprint.iacr.org/2013/552.pdf
- https://eprint.iacr.org/2015/061.pdf
- https://eprint.iacr.org/2013/491.pdf
- https://eprint.iacr.org/2019/634.pdf