

CSE 539: Applied Cryptography

Week 8: Diffie-Hellman Key Agreement

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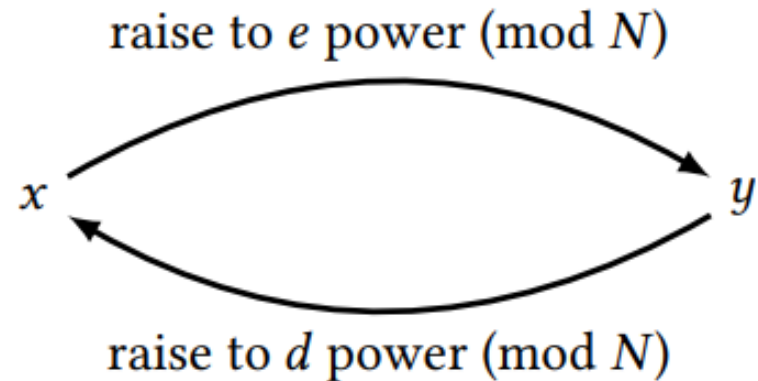
Reading: <https://joyofcryptography.com/pdf/chap14.pdf>

https://en.wikipedia.org/wiki/Diffie%E2%80%93Hellman_key_exchange

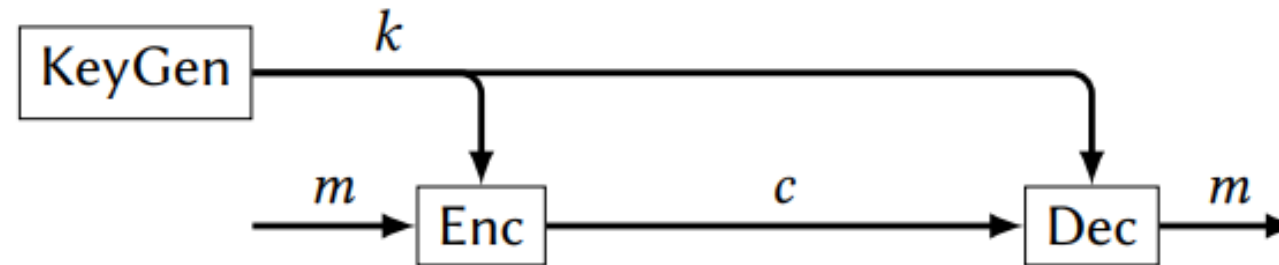
Recap: RSA

The RSA function is defined as follows:

- ▶ Let p and q be distinct primes (later we will say more about how they are chosen), and let $N = pq$. N is called the **RSA modulus**.
- ▶ Let e and d be integers such that $ed \equiv_{\phi(N)} 1$. That is, e and d are multiplicative inverses mod $\phi(N)$ – not mod N !
- ▶ The RSA function is: $x \mapsto x^e \% N$, where $x \in \mathbb{Z}_N$.
- ▶ The inverse RSA function is: $y \mapsto y^d \% N$, where $x \in \mathbb{Z}_N$.



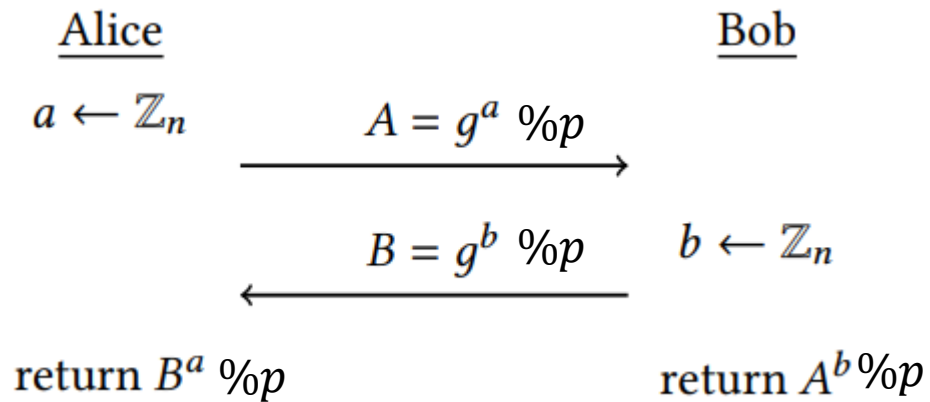
Recall: Encryption Basics & Terminology



- How to setup k ?

Diffie-Hellman Key Agreement

- Introduced in 1976
- First practical method for establishing a shared secret over an unsecured channel



Diffie-Hellman Key Agreement

- Quiz Sample: In the execution of Diffie-Hellman key agreement, Alice and Bob use the prime $p=23$ and the primitive root $g=11$. Alice chooses the secret key $a=6$. Similarly, Bob chooses the secret key $b=15$. What is the Alice & Bob's shared key?

Cyclic Groups

Definition 14.1 *Let $g \in \mathbb{Z}_n^*$. Define $\langle g \rangle_n = \{g^i \% n \mid i \in \mathbb{Z}\}$, the set of all powers of g reduced mod n . Then g is called a **generator** of $\langle g \rangle_n$, and $\langle g \rangle_n$ is called the **cyclic group generated by g mod n** .
If $\langle g \rangle_n = \mathbb{Z}_n^*$, then we say that g is a **primitive root mod n** .*

Cyclic Groups

- ▶ Any cyclic group is closed under multiplication. That is, take any $X, Y \in \mathbb{G}$; then it must be possible to write $X = g^x$ and $Y = g^y$ for some integers x, y . Using the multiplication operation of \mathbb{G} , the product is $XY = g^{x+y}$, which is also in \mathbb{G} .
- ▶ Any cyclic group is closed under inverses. Take any $X \in \mathbb{G}$; then it must be possible to write $X = g^x$ for some integer x . We can then see that $g^{-x} \in \mathbb{G}$ by definition, and $g^{-x}X = g^{-x+x} = g^0$ is the identity element. So X has a multiplicative inverse (g^{-x}) in \mathbb{G} .

Diffie-Hellman Key Agreement

Definition 14.2 *The **discrete logarithm problem** is: given $X \in \langle g \rangle$, determine a number x such that $g^x = X$.
(Discrete Log) *Here the exponentiation is with respect to the multiplication operation in $\mathbb{G} = \langle g \rangle$.**

Diffie-Hellman Key Agreement

- Quiz Sample:

Consider the following key-exchange protocol where p, g, q are public parameters

- (i) Alice chooses a random exponent $a_1 \leftarrow \mathbb{Z}_q$ and computes $h_1 = g^{a_1} \bmod p$. Alice sends h_1 to Bob
- (ii) Bob chooses two random exponents a_2, a_3 , and computes $h_2 = g^{a_2 + a_3} \bmod p$. Bob sends h_2 to Alice.
- (iii) Alice outputs a shared key $k = h_2^{a_1} \bmod p$

Show how Bob outputs the same key k ?